

Using Contracts to Guide the Search-Based Verification of Concurrent Programs

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Concurrency: necessary but challenging

- **concurrency** is necessary to exploit modern multicore architectures
- shared resources protected typically by an arbitration (**locking**) mechanism
- but this might introduce nasty **non-functional faults**
 - deadlock, livelock, starvation...
 - common symptom: lots of waiting

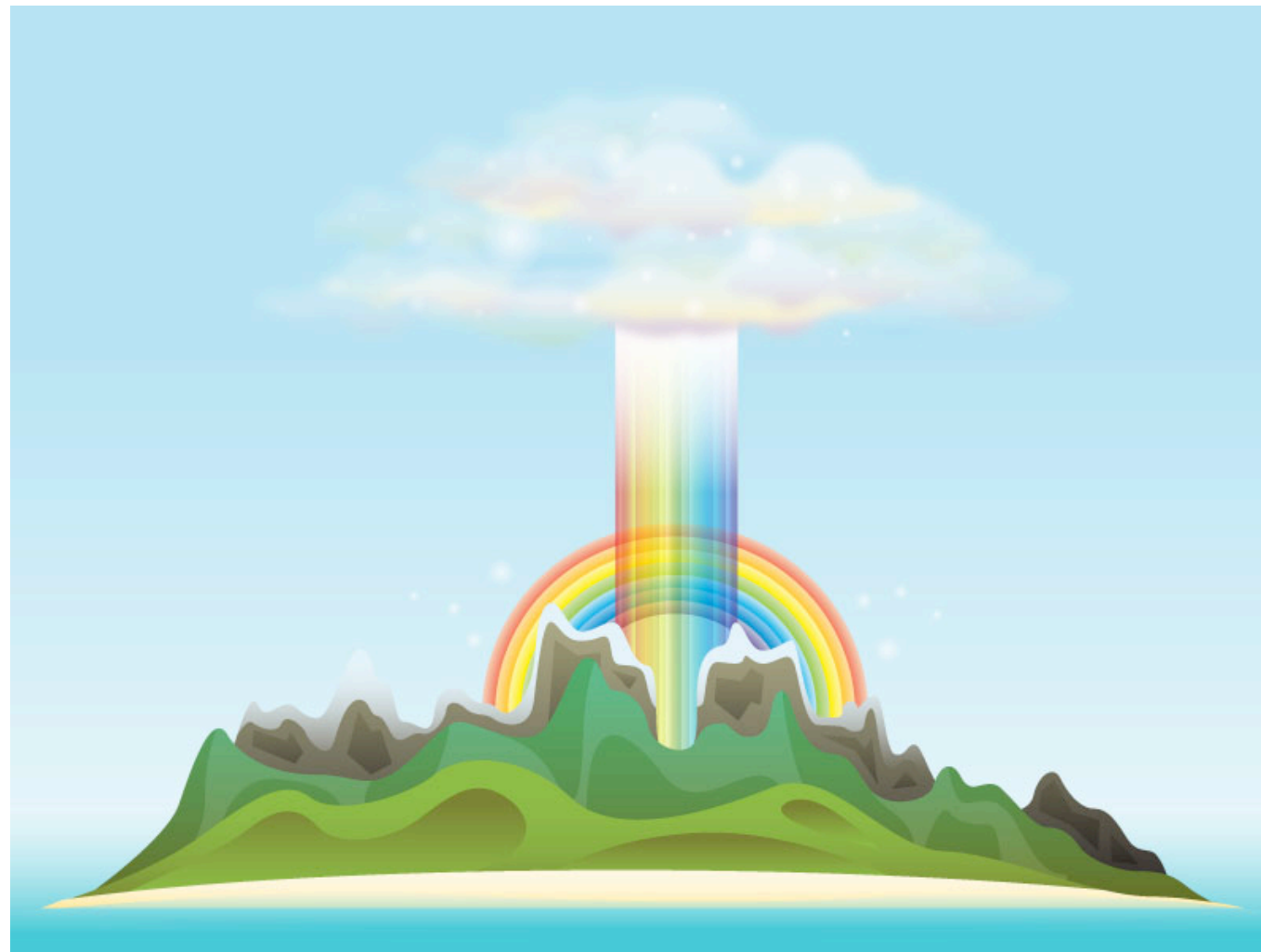
How to show the absence of such faults?

- must consider both **input data** (control flow) and **temporal orderings** of events (e.g. locking)
- exhaustive automatic techniques like **model checking** explore all possible executions



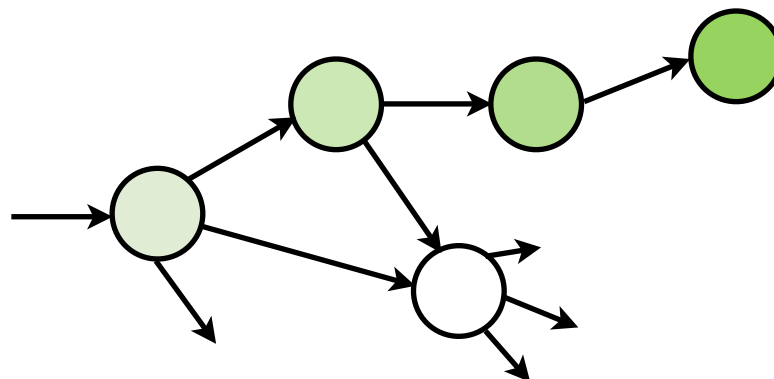
but real systems face a **combinatorial blow-up** of the state space

STATE EXPLOSION



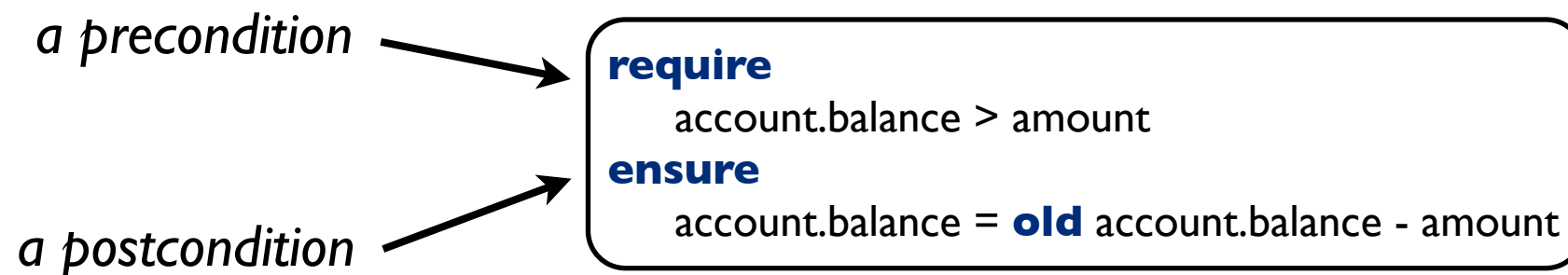
More realistic: we can search for faults (without guaranteeing we find them all)

- **metaheuristic search** can be used to evolve strategies for navigating the state space
 - *Alba et al., Godefroid & Khurshid, Shousha et al., Staunton & Clark*
- find regions more “likely” to exhibit deadlock etc.
- but **fitness functions** usually **coarse-grained**
 - maximising number of blocked threads
 - minimising number of outgoing transitions



There is an opportunity to provide more guidance

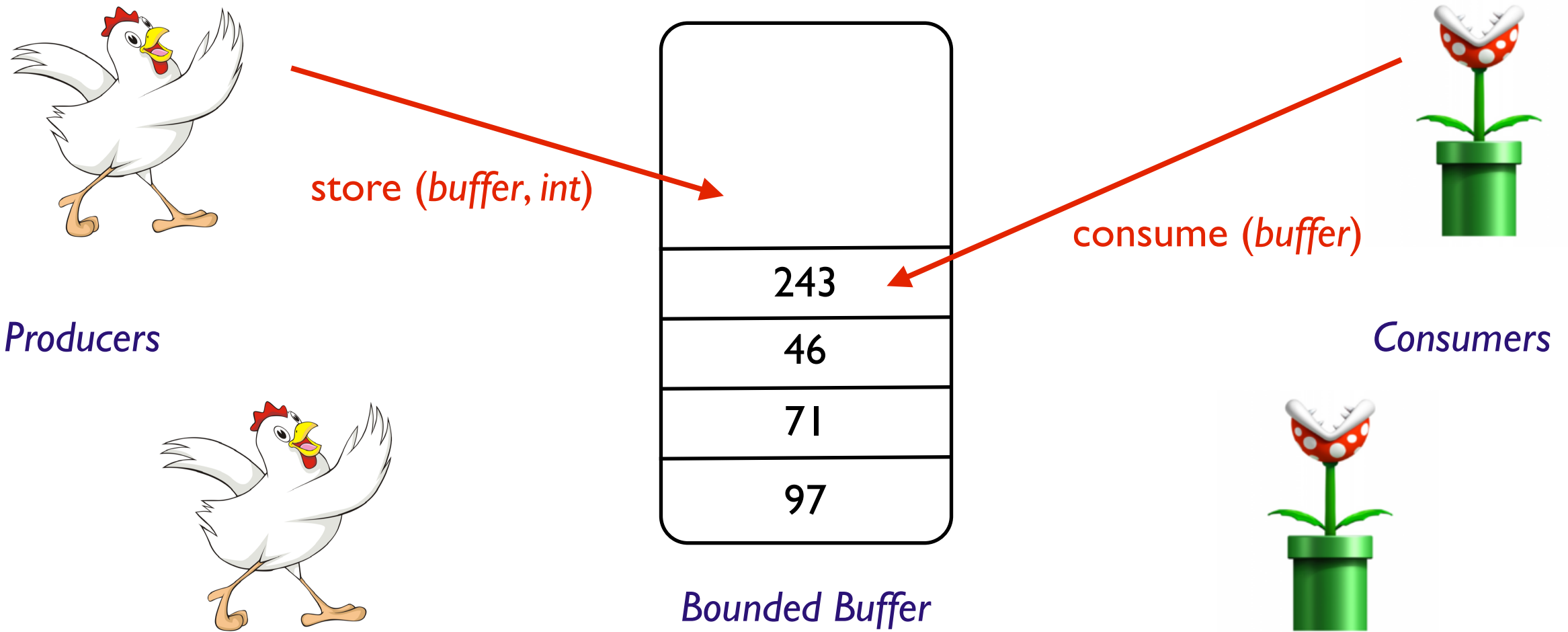
- utilise developer-specified contracts:
executable **pre- / postconditions** and **invariants**
- not full mathematical specifications; we propose to use the **Design by Contract** style (Eiffel, JML, ...)
e.g. a bank account withdrawal method



- contracts provide information that might not easily be inferred automatically

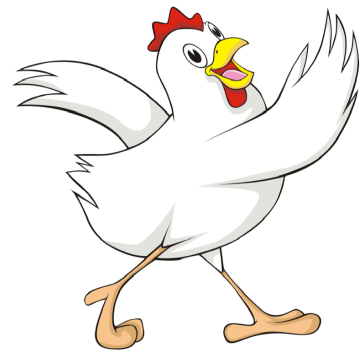
Example: producers and consumers

(SCOOP and Java versions in the paper)

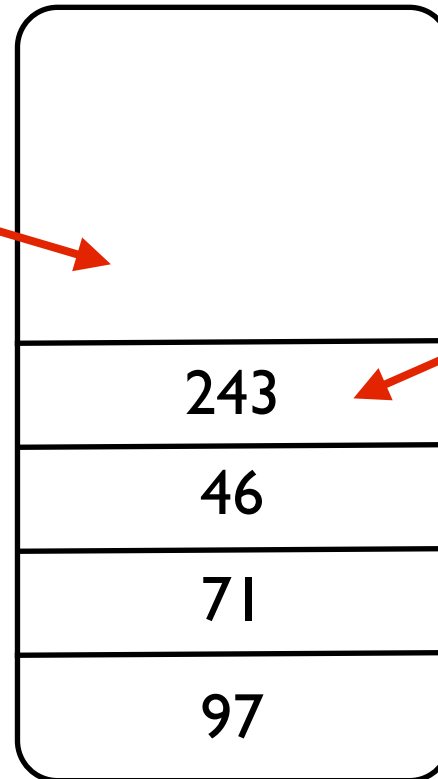


Example: producers and consumers

(SCOOP and Java versions in the paper)



store (*buffer*, *int*)



consume (*buffer*)



require
buffer.count < buffer.size

require
buffer.count > 0

Bounded Buffer

Example: producers and consumers

- the preconditions imply **regions of interest** in the state space
- program **behaviours change over such boundaries**
 - store might wait indefinitely if consumers are starved of access to full buffers
 - blocked producers may never be awoken if consumers fail to notify that the buffer is no longer full
- guide towards states that come closer to false valuations, using Korel/Tracey-like **distance metrics**



In a nutshell



We propose that additional information from **(functional) contracts** can be exploited to search for **non-functional faults** more effectively

Ongoing work; next steps

- we are working on a **prototype** using Java Pathfinder
- we are investigating how to exploit programs with **lots of contracts**
 - our example today considered two contracts in isolation
 - but how do we integrate *several* (possibly conflicting) contracts automatically and sensibly into fitness functions?
 - needs to be **self-adaptive**?
- what about **other kinds of contracts**?
 - temporal, resource usage, safety guarantees

That's all folks!



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